

# Computer Systems and Networks

ECPE 170 – Jeff Shafer – University of the Pacific

# MARIE Simulator

#### Discussion

- What was easy about programming MARIE?
- What was hard about programming MARIE?
- What tips would you share to other beginning assembly programmers?

# MARIE Programming Tip

- Writing code in assembly can be intimidating at first
  - High-level languages are (much!) easier to use
- Treat the problem like a complier would
  - Think about how to code in another language, like C++
  - Convert each construct into a series of MARIE instructions
    - ▼ Variables first, then functions, loops, if statements, ...

### SKIPCOND Operation

- - → if(AC < 0) skip next instruction
    </p>
- - if(AC == 0) skip next instruction
- - → if(AC > 0) skip next instruction
- SKIPCOND 800 = 1000 1000 0000 0000

Upper 2 bits of "address" field determine which comparison SKIPCOND will use!

### Assembly Pitfall – Program Organization

- 7 The assembler and CPU do exactly what we tell it.
  - However, that doesn't always mean it does what we intend!
- You must be very explicit when organizing your program

### Assembly Pitfall – Program Organization

Χ,

Y,

- What did the programmer intend for this code to do?
  - **7** Z=X+Y
- What does the program actually do?
  - Our variable X (0006) gets interpreted as a JNS 6 instruction!

LOADX

ADD Y

DEC 6

DEC -3

Z, DEC (

STORE Z

HALT

### Assembly Pitfall – Program Organization

- Code should always jump around any data
- Compilers did this for you in COMP 51
  - Automatically separate code and data

```
LOAD X
ADD Y
JMP SKIP
X, DEC 6
Y, DEC -3
Z DEC 0
SKIP, STORE Z
HALT
```

### Homework #9 Discussion

- Discuss algorithms for Problems 28 and 29...
- How many people submitted working .mas files that I can open in the simulator and run?
- Would you like an opportunity to re-submit a working program? (deadline: midnight?)
- Discuss schedule
  - Homework 10?
  - Quiz 3?

# Upcoming Schedule

- **尽** Monday 20<sup>th</sup> − No class
- Wednesday 22<sup>nd</sup>
  - MARIE assembly programming
  - Homework #10 Due
- Friday 24<sup>th</sup>
  - Start Chapter 5
  - **7** Quiz 3!
    - Topic: Assembly programming!
    - **↗** I will give you Table 4.7 from the book

### Subroutines



#### Subroutines

- 7 Result = addOne(input1);
- What do we need for a subroutine? (i.e. function)
  - Arguments to the function (i.e. input data)
  - Return value from the function
  - A way to jump to the function
  - A way to return from the function when finished
- Let's write a subroutine that increments a number by 1

#### Subroutine Example (Add One to Num)

```
Load Data
                    / get value
     Store Arg1 / store value as argument
     Jns AddOne / call subroutine
     Load Return / load subroutine return data
              / print it!
     Output
     Halt.
                  / terminate
Data, Dec 20
                    / original value
/ ** Subroutine **
AddOne, Dec 0
                   / return address placed here
     Load Arg1
                    / get argument
     Add One / increment it
     Store Return / save return value
     JumpI AddOne / return with value in a
                    / Empty: subroutine argument
Arg1, Dec 0
Return, Dec 0
                    / Empty: subroutine return value
```

### **Clever Tricks**



```
ORG 100
                      / Example 4.1
                      /Load address of first number to be added
     Load
           Addr
                      /Store this address is our Next pointer
     Store Next
     Load
           Num
                      /Load the number of items to be added
     Subt
           One
                      /Decrement
     Store Ctr
                      /Store this value in Ctr to control looping
Loop, Load
           Sum
                      /Load the Sum into AC
                      /Add the value pointed to by location Next
     TbbA
           Next
     Store Sum
                      /Store this sum
     Load Next
                      /Load Next
     Add
            One
                      /Increment by one to point to next address
                      /Store in our pointer Next
     Store Next
                      /Load the loop control variable
     Load
           Ctr
     Subt
                      /Subtract one from the loop control variable
           One
     Store Ctr
                      /Store this new value in loop control variable
     Skipcond 000
                      /If control variable < 0, skip next instruction
                      /Otherwise, go to Loop
     Jump
           Loop
     Halt
                      /Terminate program
Addr, Hex
            117
                      Numbers to be summed start at location 117
                      /A pointer to the next number to add
Next, Hex
                      The number of values to add
Num, Dec
Sum,
     Dec
            0
                      /The sum
                      /The loop control variable
Ctr, Hex
                      /Used to increment and decrement by 1
One, Dec
            1
            10
     Dec
                      The values to be added together
     Dec
            15
            20
     Dec
                              This is location 117
     Dec
```

- Think back to first MARIE program (Example 4.1)
- ADDR variable holds the addressof the element to be added
  - It's a pointer
  - **→** Value = 0x117

- What happens if the program changes?
  - → Say, we add a few instructions

- We would have to find the *new* starting address of the data, save its value in our assembly code, and re-run the assembler
  - Annoying!

ADDR, HEX 125

- Perfect world: The assembler lets us use a label and fills in the address from the symbol table on pass 2
  - Unfortunately we don't live in this perfect world!
  - Assembler doesn't do this directly, but we can fool it into doing something similarly useful
- Solution takes advantage of the fact that the instruction format always uses the lower 12 bits of each instruction for the address
  - i.e. this solution works for MARIE, but not necessarily other assembly languages

What happens when these changes are made to the example program?

```
100
                        Addr
               Load
101
               Store
                        Next
106
     Addi
               Next
     Addr,
111
                     List
              JnS
112
     Next,
               Hex
116
     One,
               Dec
     List,
117
               Dec
                     10
```

Recall the RTL for the ADDI instruction:

```
MAR ← X

MBR ← M[MAR]

MAR ← MBR

MBR ← M[MAR]

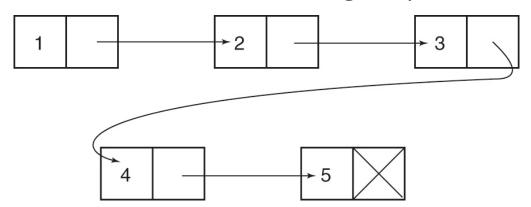
AC ← AC + MBR
```

- The third operation truncates the opcode
  - MBR = 16 bits, but MAR = 12 bits
- Result: MAR ends up with only the address of LIST

- **₹** This trick works with all instructions
  - JnS is safest since its opcode is 0
  - The resulting value placed in memory is only the 12bit address

## Homework 4.33 Tips

- Either use the JNS trick just shown, or write your program in two passes
  - Pass 1: Write the code the traverses the linked list
  - Run the assembler and look at memory addresses where it placed your program
  - Pass 2: Update the linked list memory addresses based on the assembler listing file produced



## Today's Lab

- Any or all of the following items:
  - **₹** Fix your Homework #9 and resubmit
  - → Start Homework #10
  - Assist your friends / neighbors