

Computer Systems and Networks

ECPE 170 – Jeff Shafer – University of the Pacific

Schedule

- 6 more classes left (after today)!
- Quiz 6 Tuesday, Nov 22nd
 - Input / Output (HW #15)
 - Operating Systems (HW #16)
 - Compilers & Assemblers (HW #16)
 - Review the lecture notes before the quiz (not just the homework!)
 - **Bring a Calculator**

Solid State Disks (SSDs)

See Slides from Thursday



Introduction

- Beginning Chapter 8
- System software in the form of **operating systems** and **middleware** is the glue that binds user applications and hardware together

- What software first runs when you turn on your machine?
- Not the OS! The BIOS (basic input-output system)
 - Stored on flash memory chip at known location
 - Examines system configuration
 - How many CPUs are installed?
 - How much memory is installed?
 - Where is the video card / keyboard / mouse / hard drive?
 - Assigns devices memory addresses and initializes them
 - Locates OS on disk, loads it into memory, and executes it
- A BIOS permits a single operating system to function on different computers (with different peripherals)

- The evolution of operating systems has paralleled the evolution of computer hardware
 - As hardware became more powerful, operating systems allowed people to more easily manage the power of the machine
- In the days when main memory was measured in kilobytes, and tape drives were the only form of magnetic storage, operating systems were simple resident monitor programs
 - The resident monitor could only **load**, **execute**, and **terminate programs** (command-line only!)

- In the 1960s, hardware has become powerful enough to accommodate **multiprogramming**, the concurrent execution of more than one task.
 - Multiprogramming is achieved by allocating each process a given portion of CPU time (a *timeslice*)
 - Systems were still batch oriented submit your job, and wait hours to see the results
- Interactive multiprogramming systems were called **timesharing** systems
 - You now get a *terminal* to interact with computer directly!
 - When a process is taken from the CPU and replaced by another, we say that a **context switch** has occurred

- Multiprogrammed and timesharing systems require a more complex operating system
- How to handle a context switch?
 - **♂** Save all data from current running process
 - Data includes CPU registers, page table, etc...
 - Load all data from new running process
 - ... let it run for a while ...
 - Save all data from current running process
 - Restore data from the previous running process

Operating System Tasks

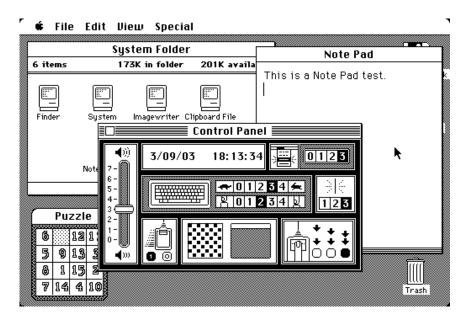


Operating System Tasks

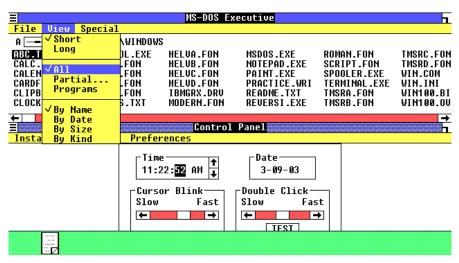
What does the OS need to do?

- Schedule processes to run
- Memory management
- Interrupt handling (manage hardware in general)
- Security (between processes)
- Network access
- Storage management (filesystem)
- Graphical user interface
 - May be a middleware layer on top of the OS

Operating systems with graphical user interfaces (GUI) were first brought to market in the 1980s



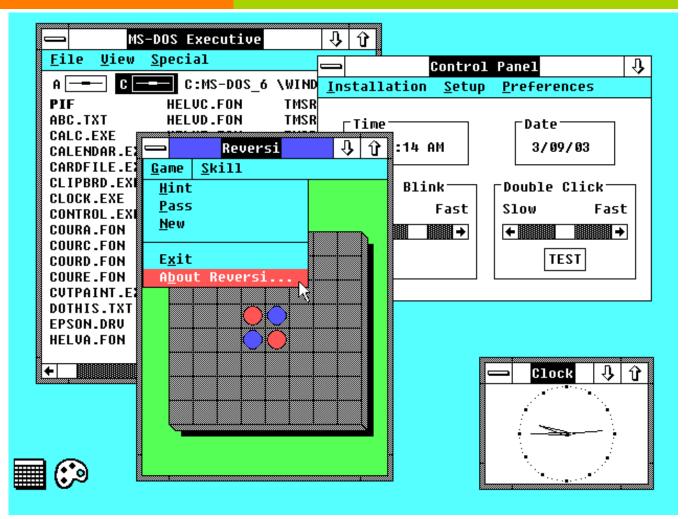
Apple Mac OS 1.0 (released 1984)



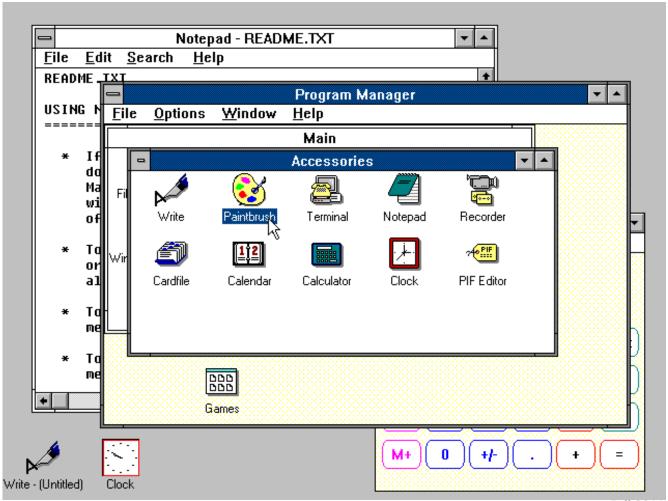
Microsoft Windows 1.0 (released 1986)

Captures from http://www.quidebookgallery.org/screenshots

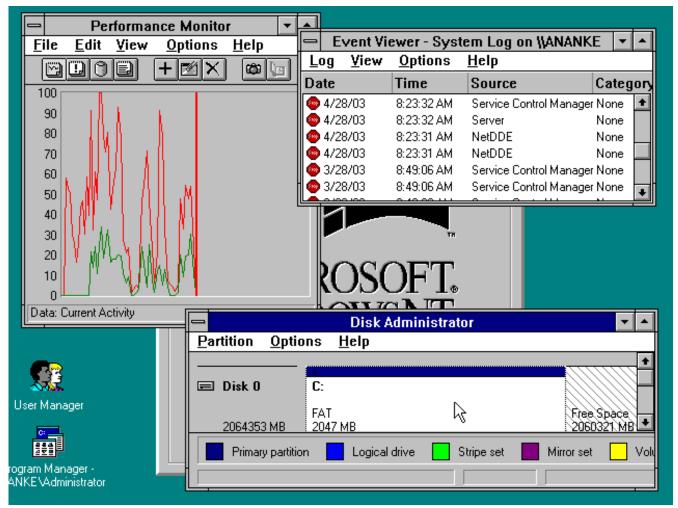
- Windows 2.0.3
- Released1987



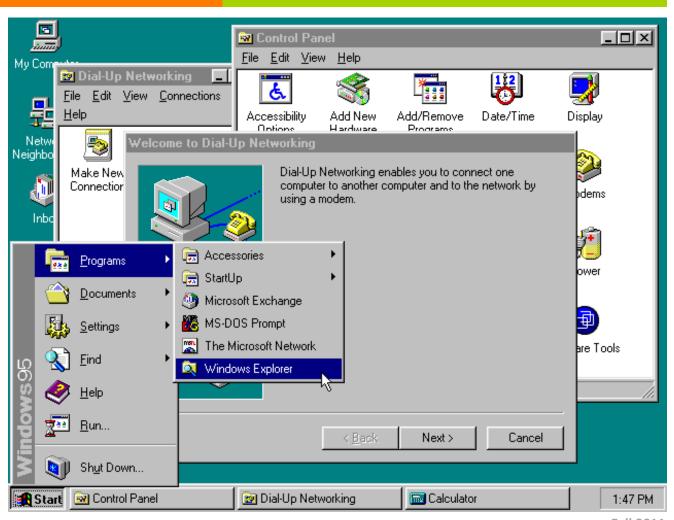
- Windows3.0
- Released1990



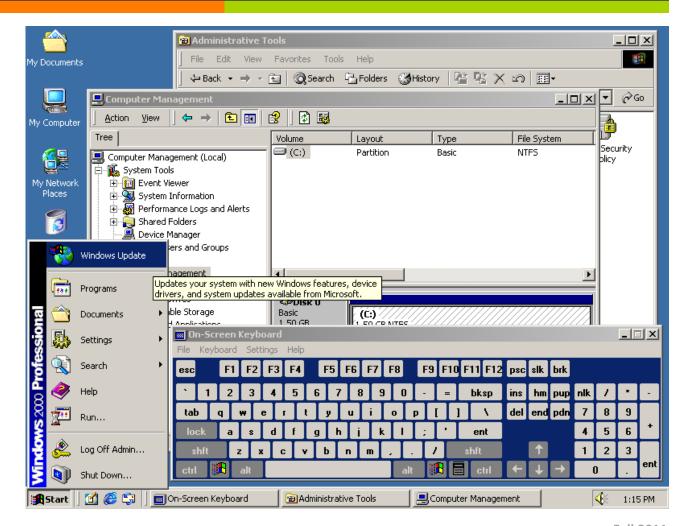
- Windows NT 3.1
- Released1993



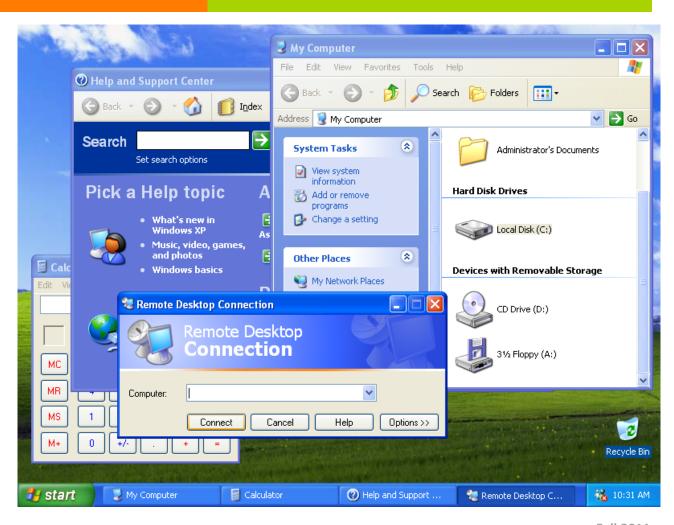
- Windows 95
- 7 Released1995



- Windows2000
- Released2000



- Windows XP
- Released2001



- From a technical perspective, the GUI is one of the least important parts of the operating system!
 - But to the users, it's the most important part
- A GUI does not even have to be part of the true OS at all
 - Windows 1.0 was just a program that ran on top of DOS, the true operating system (of that era)

Operating Systems – Processes

- **Process management** is a key operating system task
- OS must initially create processes
 - When you run your program!
- OS can allow processes to access resources
 - If resources are shared (e.g. CPU), the OS must *schedule* access to them
- OS can allow processes to **communicate** with each other
 - OS provides the mechanisms to do this
- OS must **clean up** after process finishes
 - Deallocate resources (e.g. memory, network sockets, file descriptors, etc...) that were created during process execution

Operating Systems – Scheduling

- The operating system schedules process execution
- First, the operating system determines which process shall be granted access to the CPU
 - This is long-term scheduling
- After a number of processes have been admitted, the operating system determines which one will have access to the CPU at any particular moment
 - This is short-term scheduling
- Context switches occur when a process is taken from the CPU and replaced by another process
 - CPU state (registers, current PC, etc...) is preserved during a context switch

Operating Systems – Scheduling

- Short-term scheduling can be non-preemptive or preemptive
 - Non-preemptive scheduling a process has use of the CPU until either it terminates, or must wait for resources that are temporarily unavailable
 - ▶ Preemptive scheduling each process is allocated a timeslice. When the timeslice expires, a context switch occurs
 - A context switch can also occur when a higher-priority process needs the CPU
- Which method do you think is better? (consider complexity to implement -vs- reliability)

Operating Systems – Scheduling

- Four approaches to CPU scheduling:
 - First-come, first-served Jobs are serviced in arrival sequence and run to completion if they have all of the resources they need
 - Shortest job first Smallest jobs get scheduled first. (The trouble is in knowing which jobs are shortest!)
 - **Round robin** Each job is allotted a certain amount of CPU time. A context switch occurs when the time expires
 - Priority scheduling preempts a job with a lower priority when a higher-priority job needs the CPU
- Which method(s) do modern operating systems use?

Operating Systems – Security

- Process A is forbidden from reading/modifying/writing the memory of Process B
 - **∇irtual memory** is a huge help here!
 - OS must set up virtual memory (establish page table) for each process and handle page faults when the occur
 - Wait what if I want my two programs to share memory?
- Process A has other limits besides which pages it can access
 - Ideas of other limits?
 - Amount of memory consumed
 - Number of open files on disk
 - Which files on disk can be read/written

Operating Systems – Filesystem

- OS is responsible for managing data on persistent storage
- Job of the filesystem!
 - What files exist? (i.e. names)
 - How are they organized? (i.e. paths/folders)
 - Who owns and can access them? (i.e. usernames, permissions)
 - Where are individual file blocks stored on the disk?
 - i.e. filename "database.dat" is really composed of 15823 blocks, where block 1 is located at position ...

Operating Systems – Device Management

- Manage devices
 - How do we send data to the NIC for transmission?
 - How do we render an image for display on screen?
 - How do we read a block of data from our RAID disk controller?
- Operating systems can be extended through device drivers to manage new hardware
 - Hardware vendors write software to manage their devices
 - OS provides a fixed interface (API) that driver must follow
- Common task for a device driver is **responding to interrupts** (from that device)

Operating Systems – The Kernel

- Who does all this essential work in the operating system? (besides the GUI)
 - **7** The **kernel** (i.e. the heart or core of the OS)
- Kernel performs scheduling, synchronization, memory management, interrupt handling, security and protection, etc...
- Two extremes on the design space
 - Microkernel systems Kernel provides minimal functionality and most services are carried out by external programs
 - Monolithic systems Single kernel (program) provides most of their services directly

Operating Systems – The Kernel

- Microkernel systems provide better security, easier maintenance, and portability at the expense of execution speed
 - Examples are Windows 2000, Mach, and QNX.
- Monolithic systems give faster execution speed, but are more difficult to port from one architecture to another
 - Examples are Linux, MacOS, and DOS
- Modern systems (Mac OS X, Windows XP/Vista/7) are hybrids that blend both approaches

Operating System Varients

- Several sub-categories of operating systems for special purpose applications
- Real-time Operating System
- Mobile Operating System

Real-Time Operating Systems

- Some I/O devices need to be controlled in real-time, where their correctness depends on
 - Output being <u>logically correct</u>, and
 - Output occurring at <u>correct time</u>
- Car processor must read sensors and decide whether to deploy airbags in 15-30 milliseconds
 - Similar story for traction control, ABS, etc...
 - Bad time to be running garbage collection algorithm over the memory!
 - Close enough is not good enough
- Real time operating systems (RTOS) are necessary for these types of systems

Real-Time Operating Systems

- Operating system can provide time guarantees
 - Hard real time systems have tight timing constraints
 - Soft real time systems have weaker guarantees (but are easier to create!)
- Do we want virtual memory on a RTOS?
- Do we want long-running interrupts on a RTOS? (or any interrupts?)
- Do we want garbage collection (memory management) on a RTOS?

Mobile Operating Systems

- Compare the mobile operating system on your phone (e.g. iOS, Android) versus the operating system on your laptop (e.g. Windows 7, Mac OS X)
- What does a mobile OS do that is <u>similar</u> to a desktop OS?
- What does a mobile OS do that is <u>different</u> to a desktop OS?