



Computer Systems and Networks

ECPE 170 – Jeff Shafer – University of the Pacific

Introduction to MARIE

Schedule

➤ Today

- Continue with MARIE intro
- Exam review

➤ Thursday – **Exam 1**

- Chapter 2 (Data representations)
- Chapter 3 (Digital logic)
- Part of Chapter 4 (basic organization and memory systems. Nothing on MARIE)

➤ Next Tuesday

- Meet in **CTC 115** for assembly programming
- Homework 7 due

Recap – MARIE Overview

- **How does the MARIE architecture represent positive/negative numbers?**
 - Binary, two's complement data representation
- **How is MARIE's main memory configured? (# of words, size of each word)**
 - 4K words, 16 bits wide, word-addressable

Recap – MARIE Overview

- MARIE has **seven registers** for control and data movement
 - **AC?**
 - **MAR?**
 - **MBR?**
 - **PC?**
 - **IR?**
 - **InReg?**
 - **OutReg?**

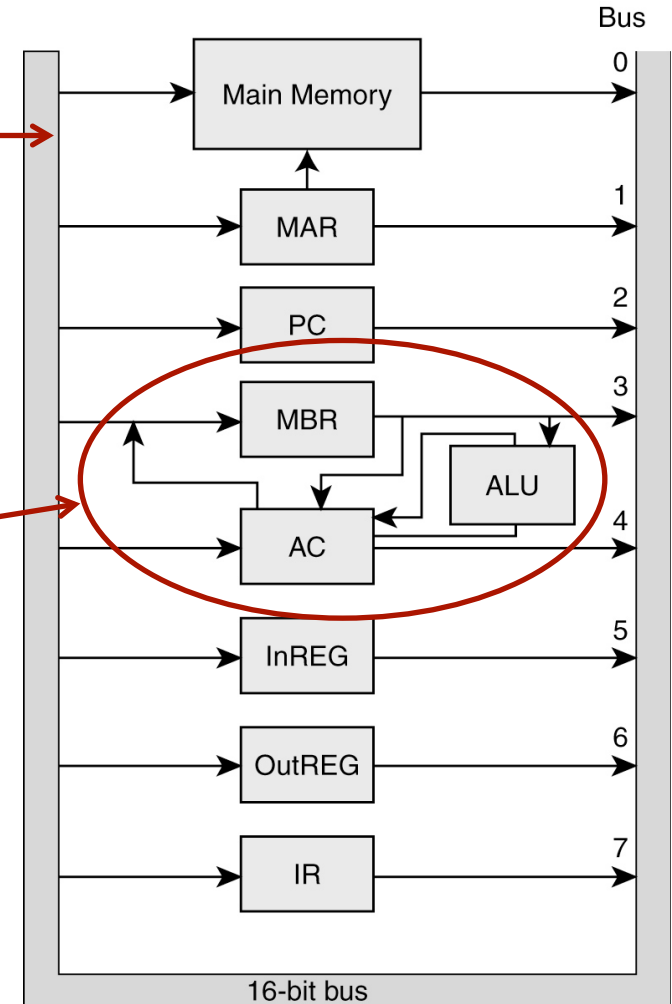
Recap – MARIE Data Path

➤ Common data bus

- Links main memory and registers
- Each device identified by unique number
- Bus has control lines that identify device used in operation

➤ Dedicated data paths

- Permits data transfer between accumulator (AC), memory buffer register (MBR), and ALU without using main data bus



Recap – MARIE ISA

- **What is an Instruction Set Architecture (ISA)?**
 - Interface between hardware and software
 - Specifies the format of processor instructions
 - Specifies the primitive operations the processor can perform

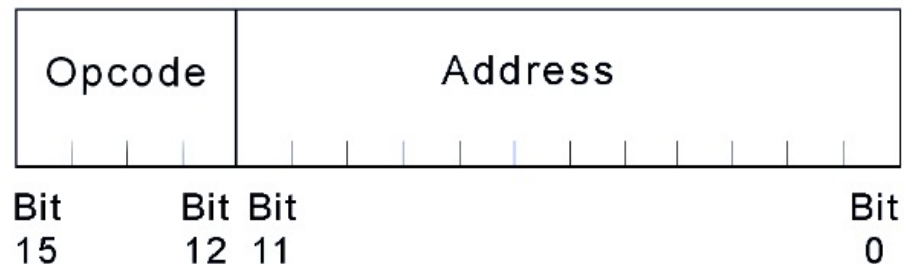
Recap – MARIE Instructions (Full)

Binary	Hex	Instruction	Meaning
0001	1	LOAD X	Load contents of address X into AC
0010	2	STORE X	Store contents of AC at address X
0011	3	ADD X	Add contents of address X to AC
0100	4	SUBT X	Subtract contents of address X from AC
0101	5	INPUT	Input value from keyboard into AC
0110	6	OUTPUT	Output value in AC to display
0111	7	HALT	Terminate program
1000	8	SKIPCOND	Skip next instruction on condition based on AC value
1001	9	JUMP X	Load value of X into PC
1010	A	CLEAR	Set AC to 0
1011	B	ADDI X	Add contents of address Mem[X] to AC
1100	C	JUMPI X	Load contents of address Mem[X] into PC
1101	D	LOADI X	Load contents of address Mem[X] into AC
1110	E	STOREI X	Store contents of AC at address Mem[X]

See table
4.7 in
book!

Recap – MARIE Instructions

- **How does MARIE format instructions in computer memory?**



- Two fields
 - **Opcode** (4 bits) – Operation code
 - **Address** (12 bits) – Address to operate to/from

Assembler



Role of Assembler

- Mnemonic instructions: LOAD 104
 - “Easy” for humans to write and understand
 - Impossible for computers to understand

- Role of **assembler**
 - Translate instructions from assembly language (*for humans*) into machine language (*for computers*)

Assembler versus Compiler

- **What's the difference between an assembler and a compiler? Which has the harder job?**
 - Assembly language → machine language
 - One-to-one correspondence
 - **Assembler** is **simple!**
 - High-level language → machine language
 - Many-to-one correspondence
 - **Compiler** is **complicated!**

Assembler Operation

- Assemblers create an **object file** (*containing machine code*) from mnemonic assembly source code in **two passes**
- Pass 1
 - Assemble as much of the program as possible
 - Builds a **symbol table** (contains memory references for all *symbols* in the program)
- Pass 2
 - Complete instructions. Fill in addresses stored in the symbol table

Assembler Operation

- Example program
 - HEX and DEC directives to specify radix of constants
- Assembler **Pass #1**
 - Create symbol table
 - Create partially-assembled instructions

Address	Instruction
100	Load X
101	Add Y
102	Store Z
103	Halt
104 X,	DEC 35
105 Y,	DEC -23
106 Z,	HEX 0000

Symbol
Table:
Name,
Address

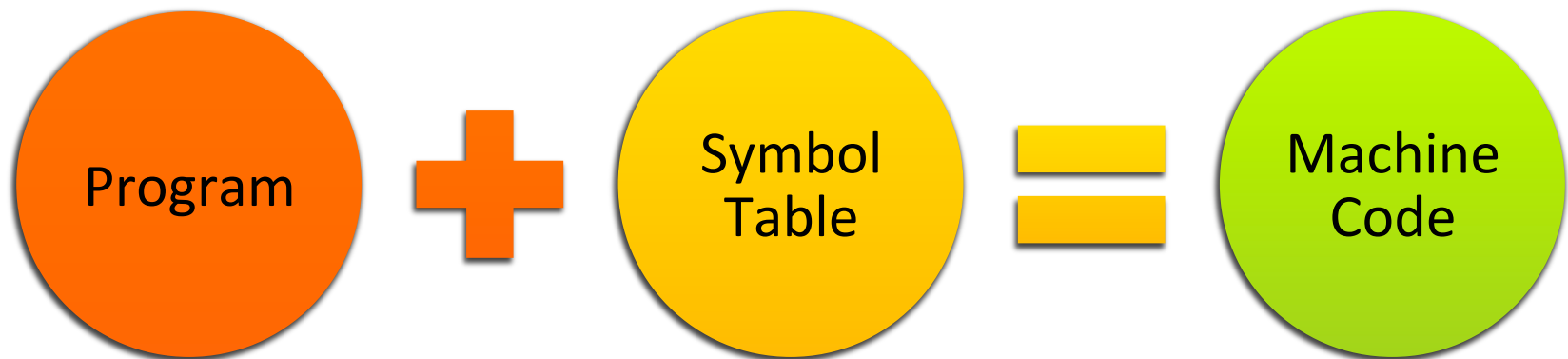
X	104
Y	105
Z	106

Partially-
Assembled
Program:

1	X
3	Y
2	Z
7	0000

Assembler Operation

- Assembler **Pass #2**
 - Fill in details from symbol table



Assembler Operation

Program:

Address	Instruction
100	Load X
101	Add Y
102	Store Z
103	Halt
104 X,	DEC 35
105 Y,	DEC -23
106 Z,	HEX 0000

Symbol Table:

X	104
Y	105
Z	106

Machine Code:

1	1	0	4
3	1	0	5
2	1	0	6
7	0	0	0
0	0	2	3
F	F	E	9
0	0	0	0

More MARIE Instructions



New Addressing Modes!

➤ **Direct addressing mode**

➤ *All the instructions covered to date...*

➤ The address of the operand is explicitly stated in the instruction

➤ **New: Indirect addressing mode**

➤ The address of the address of the operand is given in the instruction

➤ Just like **pointers** in COMP 51/53

Indirect Addressing Mode Instructions

- Four new instructions use *indirect* addressing mode: Load / store / add / jump indirect
- `LOADI X` and `STOREI X` – specified the address of the address of the operand to be loaded or stored
 - In RTL :

LOADI X

```

MAR ← X
MBR ← M[MAR]
MAR ← MBR
MBR ← M[MAR]
AC ← MBR
  
```

STOREI X

```

MAR ← X
MBR ← M[MAR]
MAR ← MBR
MBR ← AC
M[MAR] ← MBR
  
```

Indirect Addressing Mode Instructions

➤ ADDI X - Combination of LOADI X and ADD X:

➤ In RTL:

ADDI X

MAR ← X

MBR ← M[MAR]

MAR ← MBR

MBR ← M[MAR]

AC ← AC + MBR

Subroutine Instructions

- Remember subroutines? (i.e. functions)
- Machine instructions can make subroutines easier to implement
 - Jump-and-store instruction ($JNS\ X$) provides limited subroutine functionality

➤ RTL:

```
MBR ← PC
MAR ← X
M[MAR] ← MBR
MBR ← X
AC ← 1
AC ← AC + MBR
PC ← AC
```

Does JNS permit recursive calls?

No, PC is stored at address X, and we jump to address X+1. You can't do this repeatedly!

Clear Instruction

- CLEAR instruction
 - Set the contents of the accumulator to all zeroes.
- RTL for CLEAR: **AC** ← 0

MARIE Instructions (Full)

Binary	Hex	Instruction	Meaning
0001	1	LOAD X	Load contents of address X into AC
0010	2	STORE X	Store contents of AC at address X
0011	3	ADD X	Add contents of address X to AC
0100	4	SUBT X	Subtract contents of address X from AC
0101	5	INPUT	Input value from keyboard into AC
0110	6	OUTPUT	Output value in AC to display
0111	7	HALT	Terminate program
1000	8	SKIPCOND	Skip next instruction on condition based on AC value
1001	9	JUMP X	Load value of X into PC
1010	10	CLEAR	Set AC to 0
1011	11	ADDI X	Add contents of address Mem[X] to AC
1100	12	JUMPI X	Load contents of address Mem[X] into PC
1101	13	LOADI X	Load contents of address Mem[X] into AC
1110	14	STOREI X	Store contents of AC at address Mem[X]

See table
4.7 in
book!