



Computer Systems and Networks

ECPE 170 – Jeff Shafer – University of the Pacific

MIPS Assembly (Memory Fundamentals)

Lab Schedule

Activities

- **This Week**
 - Tuesday: MIPS lecture (arithmetic, branches)
 - Thursday: MIPS lecture (memory)

Assignments Due

- **Lab 10**
 - **Due by Apr 8th 5:00am**
- **Lab 11**
 - **Due by Apr 16th 5:00am**
- **Lab 12**
 - **Due by Apr 28th 5:00am**

Stub Program

```
# Declare main as a global function
.globl main

# All program code is placed after the
# .text assembler directive
.text

# The label 'main' represents the starting point
main:
    # MAIN CODE GOES HERE

# Exit the program by means of a syscall.
# There are many syscalls - pick the desired one
# by placing its code in $v0. The code for exit is "10"

li $v0, 10 # Sets $v0 to "10" to select exit syscall
syscall    # Exit

# All memory structures are placed after the
# .data assembler directive
.data

# The .word assembler directive reserves space
# in memory for a single 4-byte word (or multiple 4-byte words)
# and assigns that memory location an initial value
# (or a comma separated list of initial values)
#For example:
#value: .word 12
```

MIPS Memory Access



Memory

- Challenge: **Limited supply of registers**
 - Physical limitation: We can't put more on the processor chip, and maintain their current speed
 - *Many elements compete for space in the CPU...*
- Solution: **Store data in memory**
- MIPS provides instructions that transfer data between memory and registers

MIPS Memory Declarations

- All of the memory values must be declared in the `.data` section of the code
 - You ask the assembler to reserve a region of memory in the *data* section and refer to that region with a *label*

- Examples
 - Declare a 32-bit word with initial value of 12:
`Z: .word 12`
 - Declare a 256 byte region of memory (could be 64 integers, 256 chars, etc...)
`array: .space 256`
 - Declare a null-terminated string with initial value
`msg: .asciiz "Hello world!"`

Memory Fundamentals

MIPS cannot directly manipulate data in memory!

Data must be moved to a register first! (And results must be saved to a register when finished)

This is a common design in *RISC-style* machines: a *load-store* architecture

Memory Fundamentals

Yes, it's a **pain** to keep moving data between registers and memory.

But consider it your *motivation* to reduce the number of memory accesses. That will **improve program performance!**

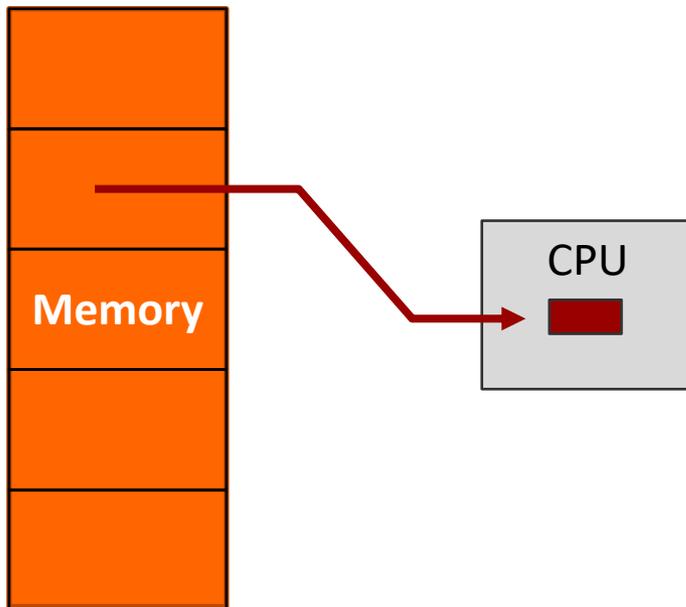
Memory Fundamentals

- Four questions to ask when accessing memory:
1. What **direction** do I want to copy data?
(i.e. to memory, or from memory?)
 2. What is the specific **memory address**?
 3. What is the specific **register name**? (or number)
 4. How **much data** do I want to move?

Memory – Fundamental Operations

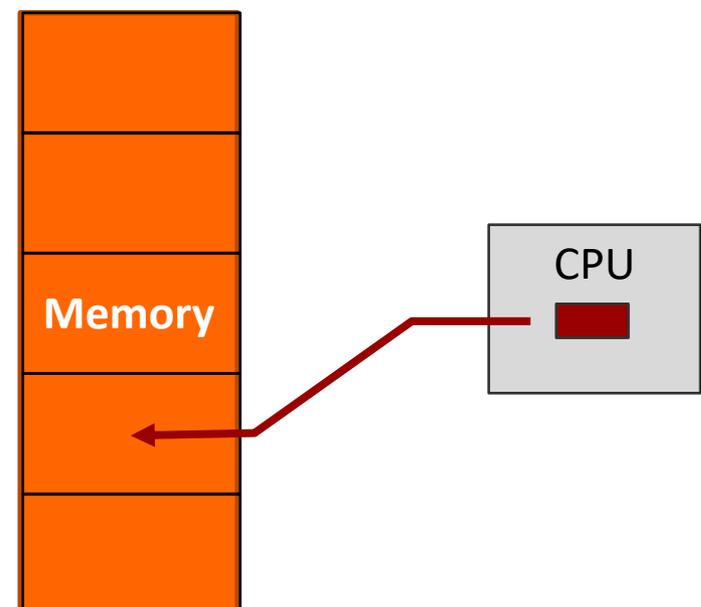
Load

- Copy data from memory to register



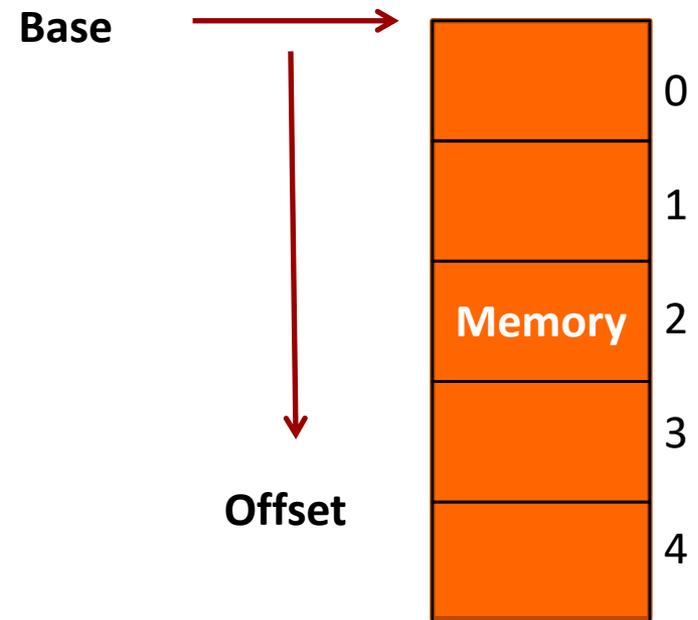
Store

- Copy data from register to memory



Memory – Determining Address

- There are many ways to calculate the desired memory address
 - These are called *addressing modes*
 - We'll just learn one mode now:
base + offset
- The base address could be HUGE! (32 bits)
 - We'll place it in a **register**
- The offset is typically small
 - We'll directly include it in the instruction as an "immediate"



MIPS notation: `offset(base)`

Memory – Register Name

- What is the name of the register to use as either the data destination (for a *load*) or a data source (for a *store*)?
- Use the same register names previously learned

Memory - Data Transfer Size

- How much data do I want to load or store?
 - A full word? **(32 bits)**
 - A “half word”? **(16 bits)**
 - A byte? **(8 bits)**
- **We’ll have a different instruction for each quantity of data**
- No option to load an entire array!
 - Will need a loop that loads 1 element at a time...

Memory – Data Transfer Instructions

➤ Load (copy from memory to register)

Word: `lw <reg>, <offset>(<base addr reg>)`

Byte: `lb <reg>, <offset>(<base addr reg>)`

➤ Store (copy from register to memory)

Word: `sw <reg>, <offset>(<base addr reg>)`

Byte: `sb <reg>, <offset>(<base addr reg>)`

Register

Memory Location

Example

➤ What will this instruction do?

```
lw $s1, 20($s2)
```

- Load word copies from memory to register:
 - Base address: stored in register \$s2
 - Offset: 20 bytes
 - Destination register: \$s1
 - Amount of data transferred: 1 word (32 bits)

Problem 1: Simple Program

- Declare memory variables *A*, *B*, and *C*, initialized to 20, 45, and 0, respectively. In *main*, set *C* to sum of *A* and *B*.

```
.globl main
.text
main: #Main goes here

    li $v0, 10 #v0 argument set to 10
                    # for system call "exit"
    syscall
.data             #Data goes in this section
```

Aside – Compiler

- **When programming in C / C++, are your variables (int, float, char, ...) stored in memory or in registers?**
- **Answer: It depends**
- **Compiler will choose** where to place variables
 - Registers: Loop counters, frequently accessed scalar values, variables local to a procedure
 - Memory: Arrays, infrequently accessed data values

MIPS Array Access



Arrays Revisited

- Name of the array is the address of first element

```
int array[20];  
printf("Address of first element:%u",array);
```

- Values are spaced by the size of the data

- Integers – Spaced by 4 bytes

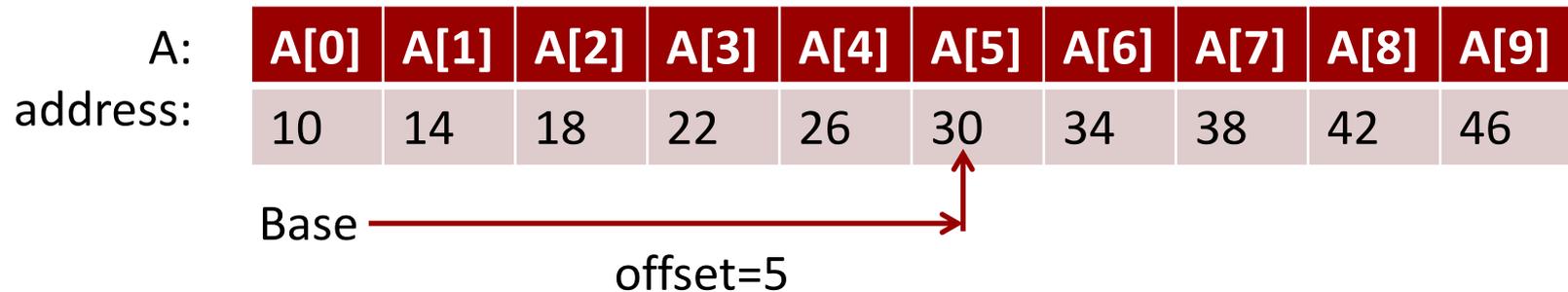
- Doubles – Spaced by 8 bytes

```
int array[20];  
printf("Address of the first element:%u",  
      &array[0]); // Say it prints 65530  
printf("Address of the second element:%u",  
      &array[1]); // Will print 65534
```

Arrays Revisited

Base offset addressing / Indexed Addressing:

`A[5]`, `array[i]`, ...



Pointer arithmetic:

```
int array[10];
```

```
printf("array[5]:%u", *(array+5));
```

//Adds **20 bytes** to base address to access `array[5]`

Remember, in C, pointer arithmetic is done with respect to data size!

Problem 2: Arrays Revisited

- Write a C for-loop to print the values of a 1-D array of size N using:
1. Indexed addressing
 2. Pointer arithmetic

Task : Write Code

➔ **Write MIPS assembly for:**

$g = h + \text{array}[16]$

(Array of words. Can leave g and h in registers)

Map:

$\$s1 = g$
 $\$s2 = h$
 $\$s3 = \text{base}$
address of
array

Code:

```
# Assume $s3 is already set
lw $t0, 16($s3)
add $s1, $s2, $t0
```

Memory Address

➤ Slight flaw in previous solution

- The programmer intended to load the 16th array element
- Each element is 4 bytes (1 word)
- The offset is in bytes
- $16 * 4 = 64$

Correct Code:

```
# Assume $s3 is already set  
lw $t0, 64($s3)  
add $s1, $s2, $t0
```

C vs. MIPS

C Programming

- ➔ C has the format:
`base [offset]`
- ➔ The C compiler multiplies the `offset` with the size of the data to compute the correct offset in bytes

MIPS Programming

- ➔ MIPS has the format:
`offset (<base-addr-reg>)`
- ➔ In MIPS, **YOU** multiply the offset with size of the data to compute the correct `offset` in bytes

Problem 3: Base Offset Addressing

➔ **Write MIPS assembly for:**

array[12] = h + array[8]

(Array of words. Assume h is in register)

Map:

\$s2 = h
\$s3 = base
address of
array
\$t1 = temp

Code:

```
# Assume $s3 is already set  
lw $t0, 32($s3)  
add $t1, $s2, $t0  
sw $t1, 48($s3)
```

Problem 4: Pointer Arithmetic

➔ **Write MIPS assembly for:**

$$g = h + \text{array}[i]$$

(Array of words. Assume g , h , and i are in registers)

Map:

```
$s1 = g  
$s2 = h  
$s3 = base  
address of  
array  
$s4 = i
```

Code:

```
# "Multiply" i by 4  
add $t1, $s4, $s4    # x2  
add $t1, $t1, $t1    # x2 again  
# Get addr of array[i]  
add $t1, $t1, $s3  
# Load array[i]  
lw $t0, 0($t1)  
# Compute add  
add $s1, $s2, $t0
```

P4

Addresses

- Tip: To get the address of a label in the `.data` section, use the “load address” instructions (`la`)

```
la <reg>, label
```

Example:

```
# Load the starting address of  
# the label 'array' into $s0  
la $s0, array
```

Problem 5: Full Program

- Write a **complete MIPS program** which implements the C code below. Test your program in QtSPIM.

```
int array[7]; // Store in memory
int main()
{
    int i=0; // Store in register
    array[0]=5;
    array[1]=4;
    for(i=2; i<7; i++)
        array[i] = array[i-2] + array[i-1];
}
```